

MODULE 8

MODULARITY

Course Material for Modularity

Chapter 11

P = primair, I = Illustratie, O = overslaan

P	11.1	Introduction	560
P	11.2	Datapaths in Digital Processor Architectures	560-561
P	11.3	The Adder	561
P	11.3.1	The Binary Adder: Definitions	561-564
P	11.3.2	The Full Adder: Circuit Design Consideration	564-578
I	11.3.2	Manchester Carry Chain Adder	568-570
O	11.3.3	The Binary Adder: Logic Design Considerations	571-586
P	11.4	The Multiplier	586
P	11.4.1	The Multiplier: Definitions	586-587
P	11.4.2	Partial-Product Generation	587-589
P	11.4.3	Partial-Product Accumulation	589-592
O		The Tree Multiplier	592-593
I	11.4.4	Final Addition	593-594
P	11.5	The Shifter	595
P	11.5.1	Barrel Shifter	595-596
P	11.5.2	Logarithmic Shifter	596
O	11.6	Other Arithmetic Operators	596-600
O	11.7	Power and speed trade-offs	600-618
P	11.8	Perspective: Design as a Trade-off	618-619
P	11.9	Summary	619-620

Outline

- **Background on Modular Design**
 - **Hierarchy, reuse, regularity**
 - **Architecture, bit-slicing**
- **Adder Design**
- **Multiplier Design**
- **Shifter Design**
- **Layout Strategies (regularity)**
- **Design as a Trade-Off**

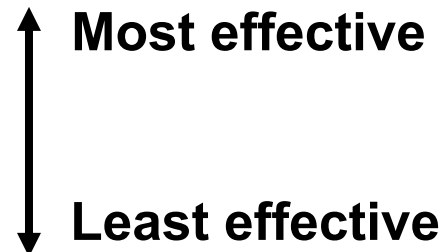
**contains a lot
of reminders**

**Get further appreciation of some
system level design issues**

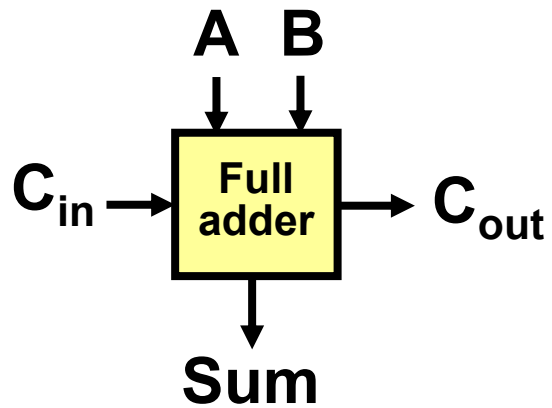
Adder Design

- Adders are fundamental building blocks
 - Digital filtering (DSP): MP3 en/decoder, GSM, GPS, ...
 - Data processing
 - Multiplication
 - Address arithmetic
 - ...
- Good performance is key
- Many architectures
 - ✓ ■ Static adder
 - ✗ ■ Dynamic adder (Manchester Carry Chain)
 - ✗ ■ Pipelined Adder
 - ✗ ■ Carry-Bypass, Carry Lookahead, Carry Select
 - ✗ ■ ...
- Design trade-offs, optimization
 - ✓ ■ Architecture level
 - ✓ ■ Logic level
 - ✓ ■ Circuit level
 - ✓ ■ Layout level

Also see Digital Systems – Katz § 5.5



Full-Adder



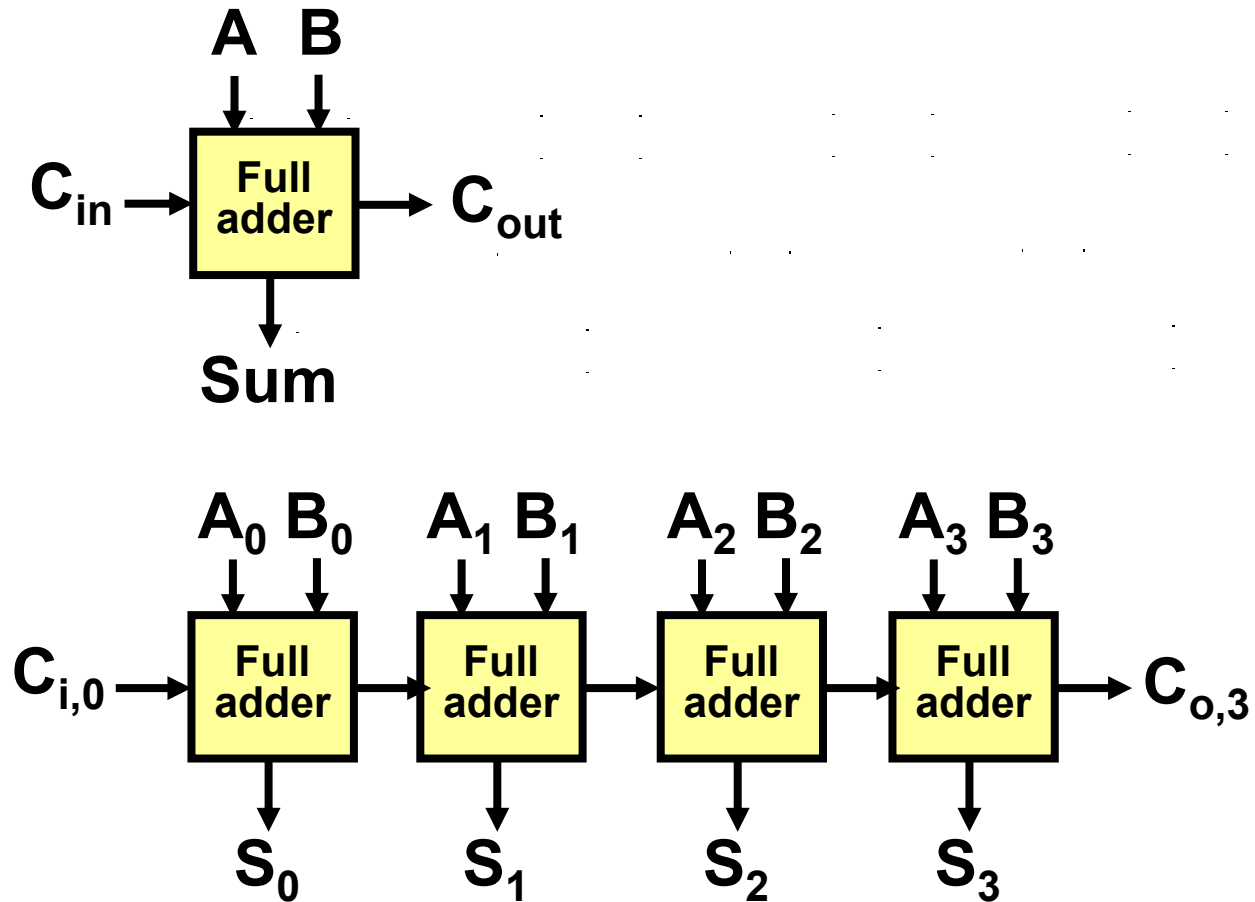
Add three one-bit numbers

Equivalently: count # 1's in A, B, C_i

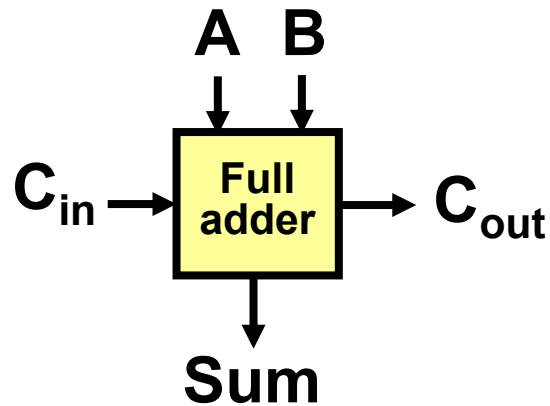
Output as 2-bit number $\langle C_o S \rangle$

A	B	C_i	S	C_o
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

The Ripple-Carry Adder



The Binary Adder



	\bar{c}	c
$\bar{A}\bar{B}$		1
$\bar{A}B$	1	
AB		1
$A\bar{B}$	1	

(a) SUM

	\bar{c}	c
$\bar{A}\bar{B}$		
$\bar{A}B$		1
AB	1	1
$A\bar{B}$		1

(b) CARRY

$AB + BC + AC$

$$\begin{aligned}
 S &= A \oplus B \oplus C_i \\
 &= \bar{A}\bar{B}C_i + \bar{A}B\bar{C}_i + A\bar{B}\bar{C}_i + ABC_i \\
 C_0 &= AB + BC_i + AC_i
 \end{aligned}$$

AND-OR expressions for sum and carry

Naive Complementary CMOS Implementation

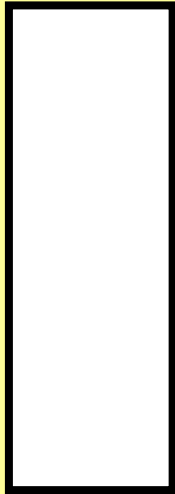
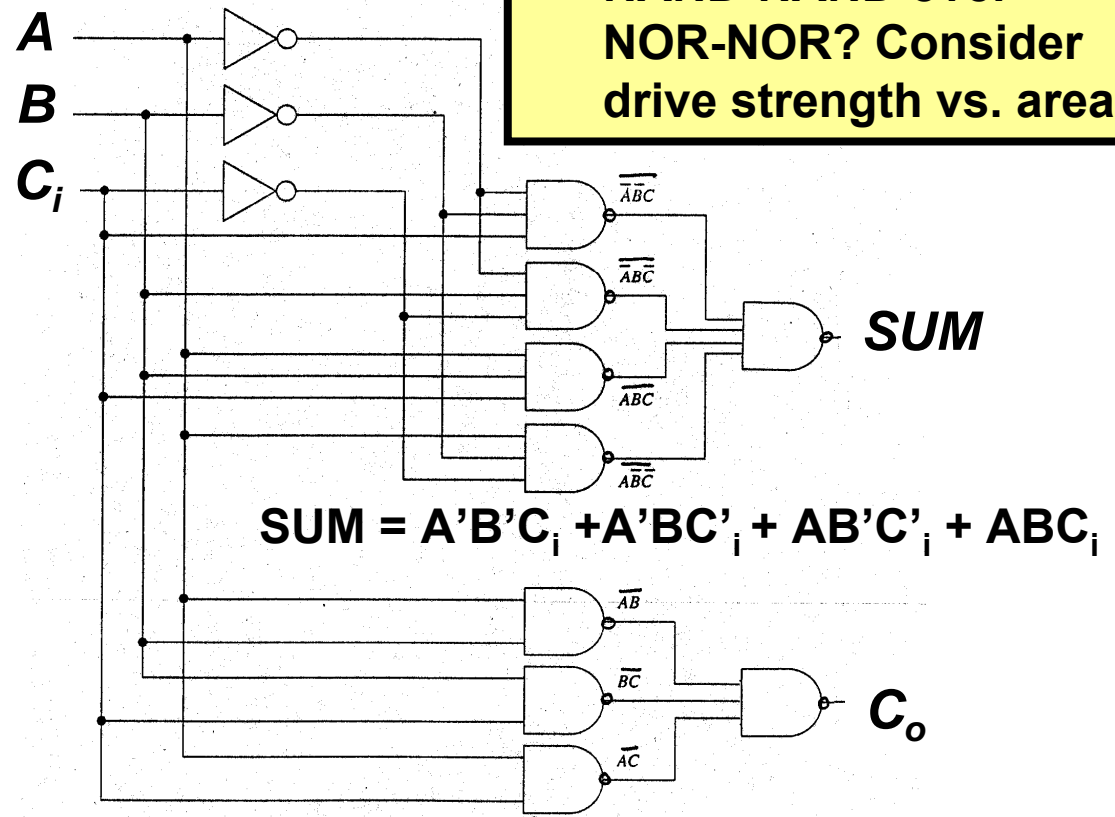
- Use DeMorgan to convert AND-OR expressions for *SUM* and *CARRY* to NAND-NAND

- $PQ + RS = \overline{\overline{PQ} \overline{RS}}$ (example)

Q: What is advantage of NAND-NAND over NOR-NOR? Consider drive strength vs. area

Transistor Count

3 × INVERT
 3 × NAND-2
 5 × NAND-3
 1 × NAND-4

Can do better using more clever boolean factoring, but...

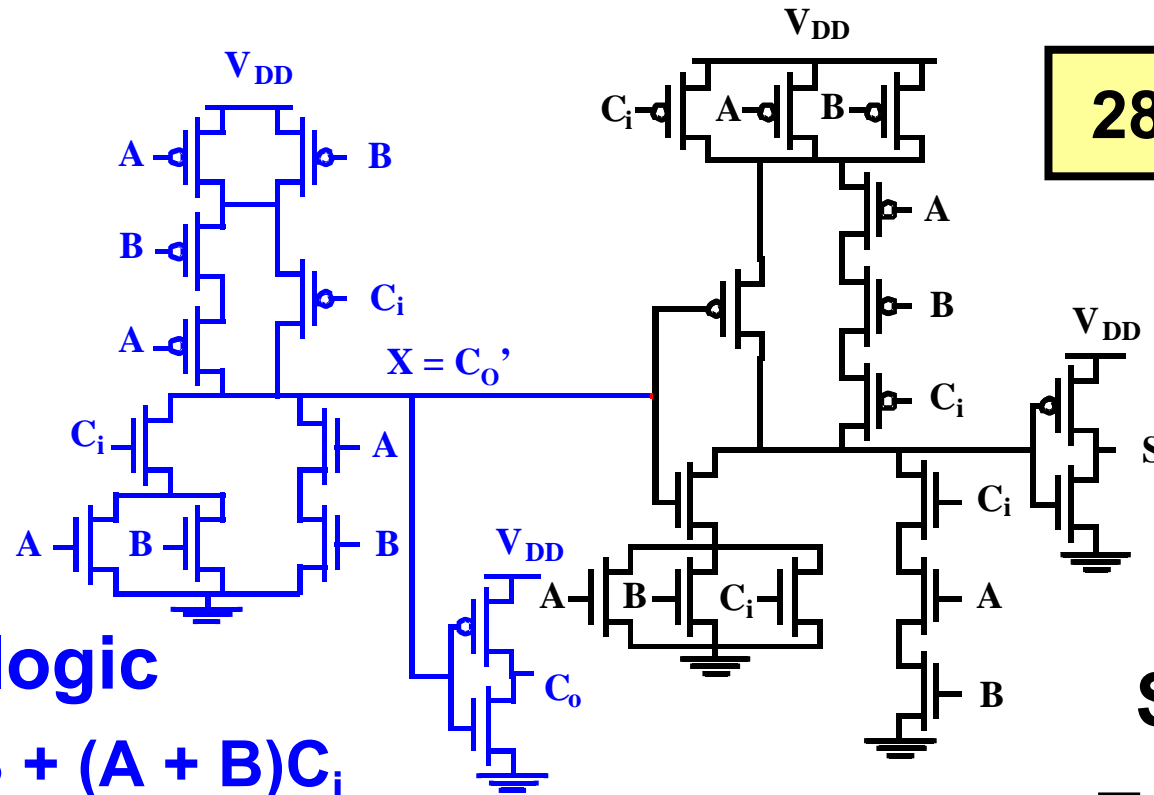
Full-Adder Boolean Factoring

$$\begin{aligned} S &= \bar{A}\bar{B}\bar{C}_i + \bar{A}B\bar{C}_i + A\bar{B}\bar{C}_i + ABC_i \\ &= ABC_i + \bar{C}_0(A + B + C_i) \end{aligned}$$

$$\begin{aligned} C_o &= AB + BC_i + AC_i \\ &= AB + (A + B)C_i \end{aligned}$$

A	B	C_i	S	C_o
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

Improved Complementary Static Full Adder



28 Transistors

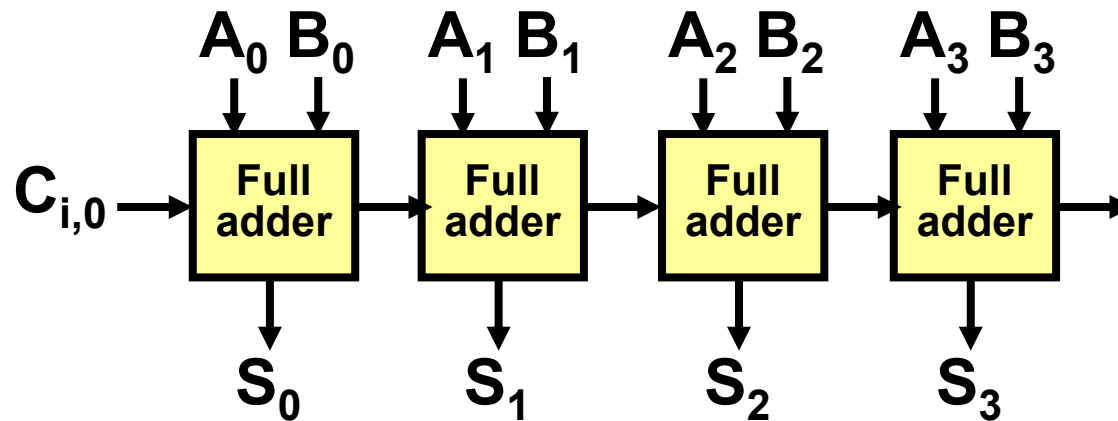
Carry logic

$$C_0 = AB + (A + B)C_i$$

Sum logic

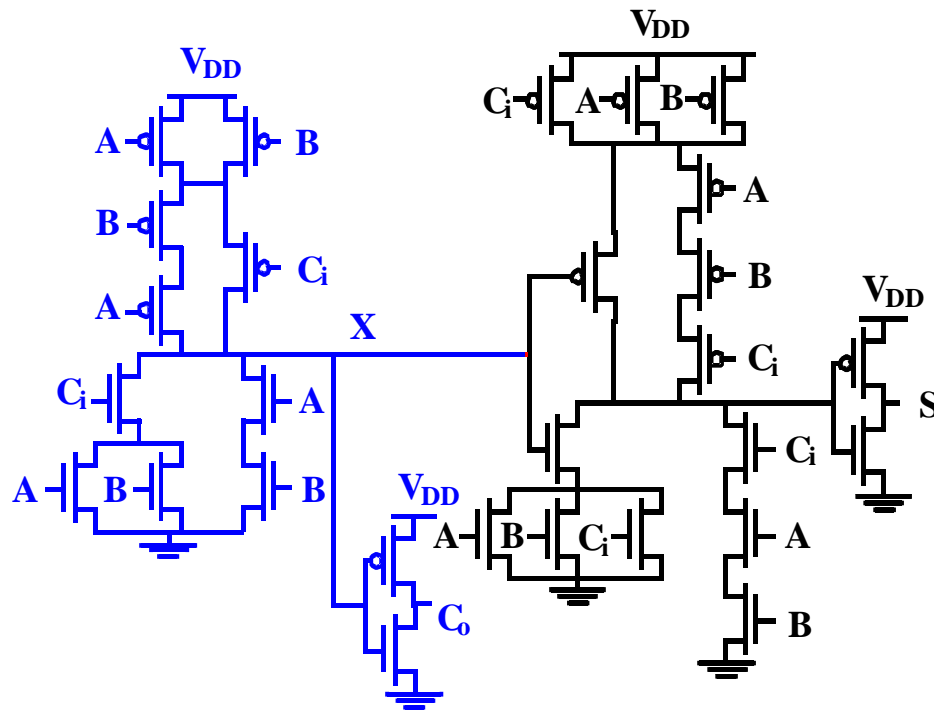
$$S = ABC_i + \bar{C}_0(A + B + C_i)$$

Ripple-Carry Adder Delay



- Worst case delay through full carry path (ripple carry)
- Linear with the number of bits (N)
- $T_{\text{adder}} = (N-1) T_{\text{carry}} + \text{Max} (T_{\text{carry}}, T_{\text{sum}})$
- $T_{\text{adder}} = O(N)$ “ T_{adder} is of Order N” *means linear with N*
- **Goal: Make the fastest possible carry path circuit**

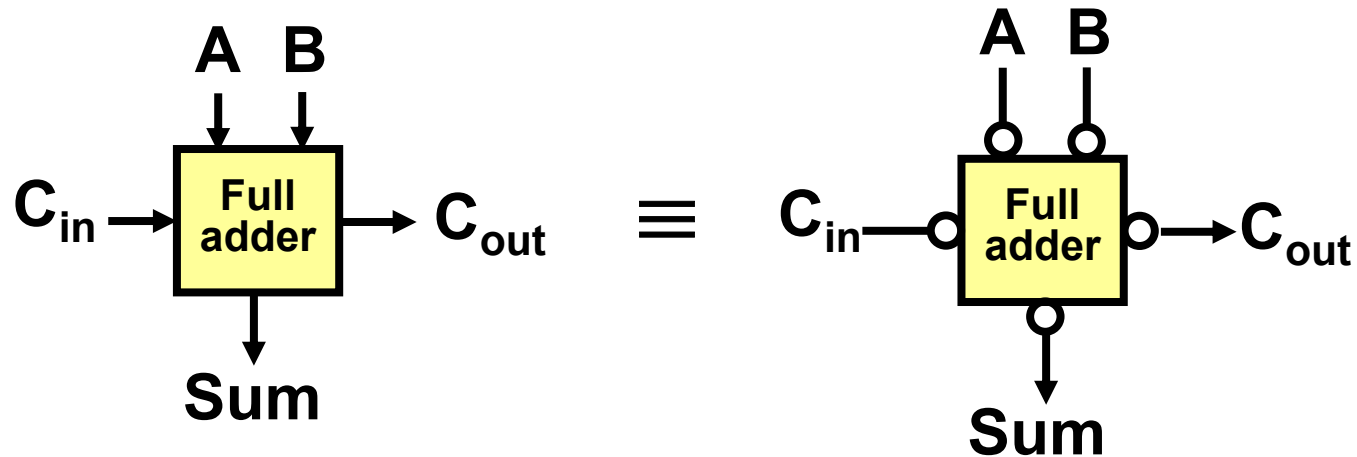
Adder Evaluation



Carry Chain:

- Long PMOS chains
- High C at X
- 2 (inverting) stages

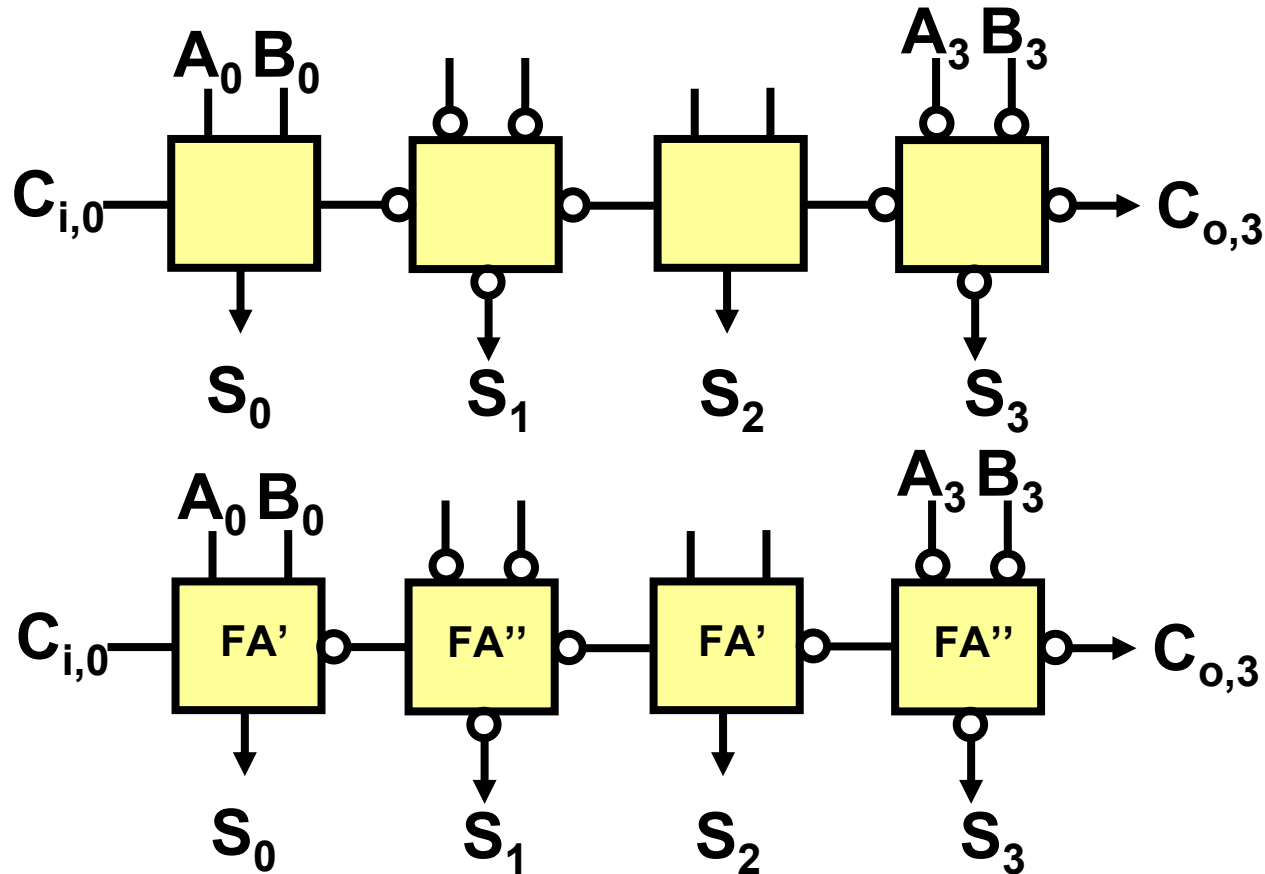
Inversion Property



$$\bar{S}(A, B, C_i) = S(\bar{A}, \bar{B}, \bar{C}_i)$$

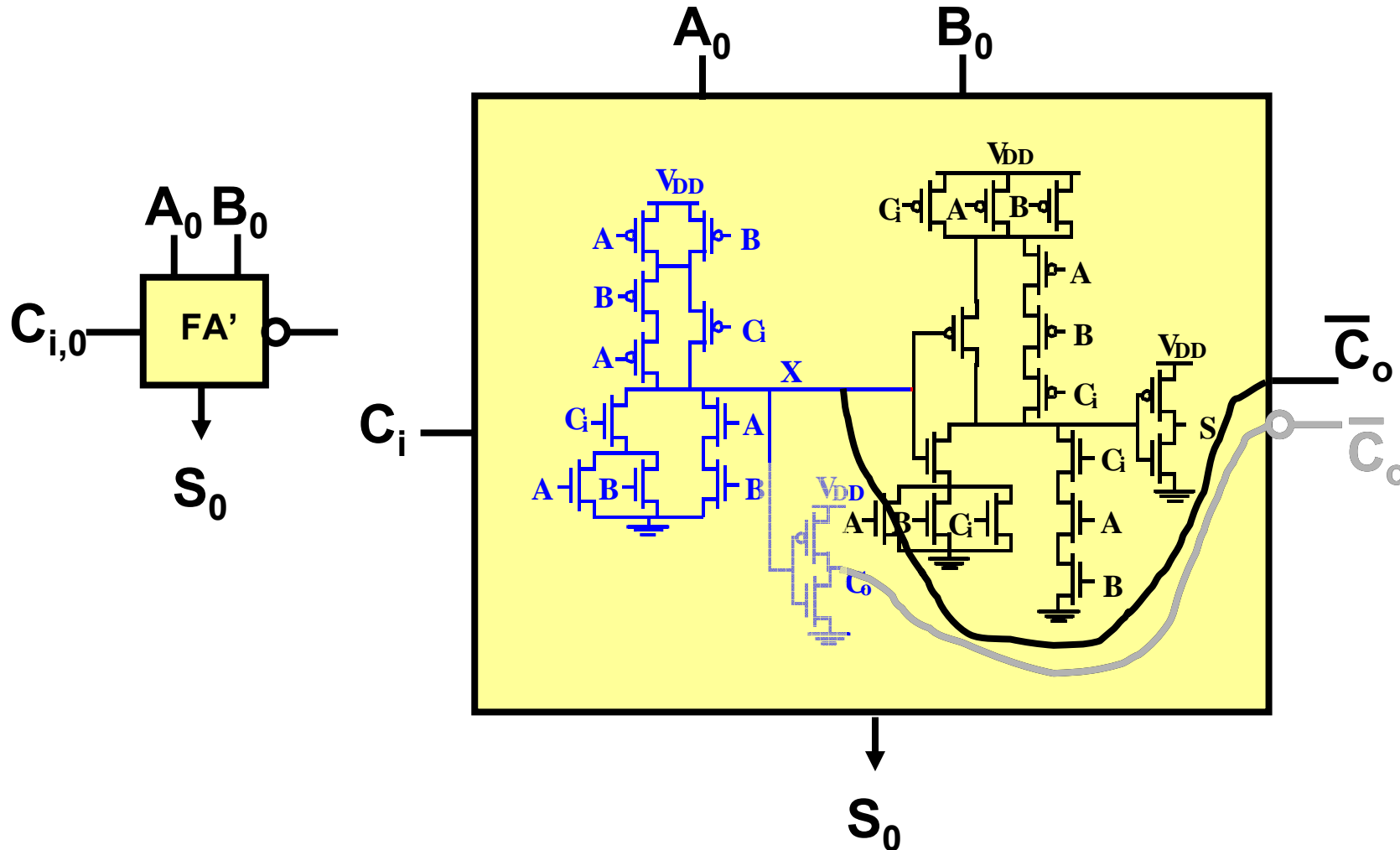
$$\bar{C}_0(A, B, C_i) = C_0(\bar{A}, \bar{B}, \bar{C}_i)$$

Minimize Critical Path by Reducing Inverting Stages



- Can eliminate inverter in carry from each FA
- Need 2 different types of cells, but both with inverting carry – will require only one stage per bit

Eliminate Inverter In Carry



Multiplier Design

- **Multipliers are fundamental building blocks *too***
 - **Digital filtering (DSP): MP3 en/decoder, GSM, GPS, ...**
 - **Data processing**
 - **Address arithmetic**
 - **...**
- **Good performance is key, often they are the performance bottleneck**
- **Multipliers are complex *arrays of adders***
- **Many architectures**
 - ✓ ■ **Basic Array Multiplier**
 - ✗ ■ **Bit-serial**
 - ✗ ■ **Booth-encoding multiplier**
 - ✗ ■ **Baugh-Wooley multiplier**
 - ✗ ■ **Wallace tree multiplier**
 - ✗ ■ **...**
- **Design trade-offs, optimization**
 - ✓ ■ **Architecture level, Logic level, Circuit level, Layout level**

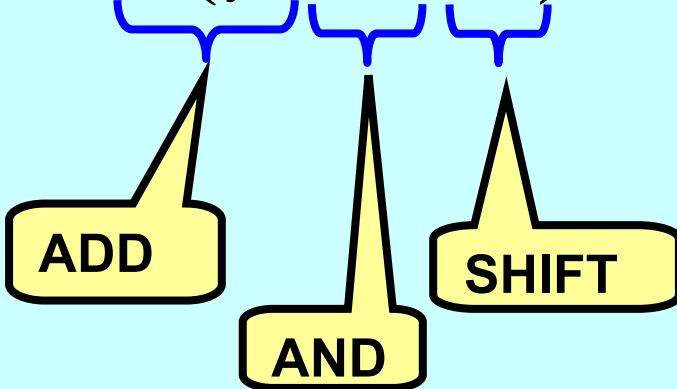
The Binary Multiplication

$$X = \sum_{i=0}^{M-1} X_i 2^i \quad Y = \sum_{j=0}^{N-1} Y_j 2^j$$

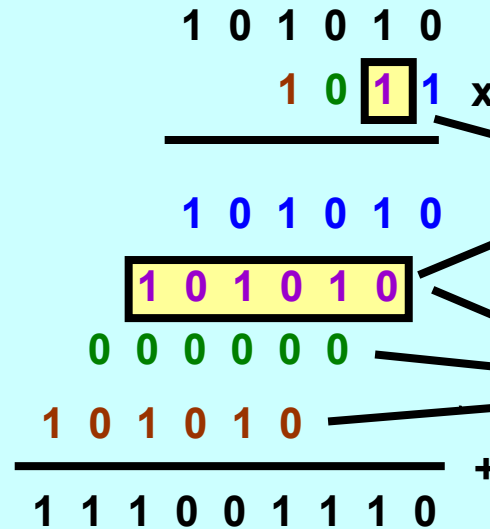
$$Z = X \times Y = \sum_{k=0}^{M+N-1} Z_k 2^k$$

$$= \left(\sum_{i=0}^{M-1} X_i 2^i \right) \left(\sum_{j=0}^{N-1} Y_j 2^j \right)$$

$$= \sum_{i=0}^{M-1} \left(\sum_{j=0}^{N-1} X_i Y_j 2^{i+j} \right)$$



Example: 42 x 11 = 462



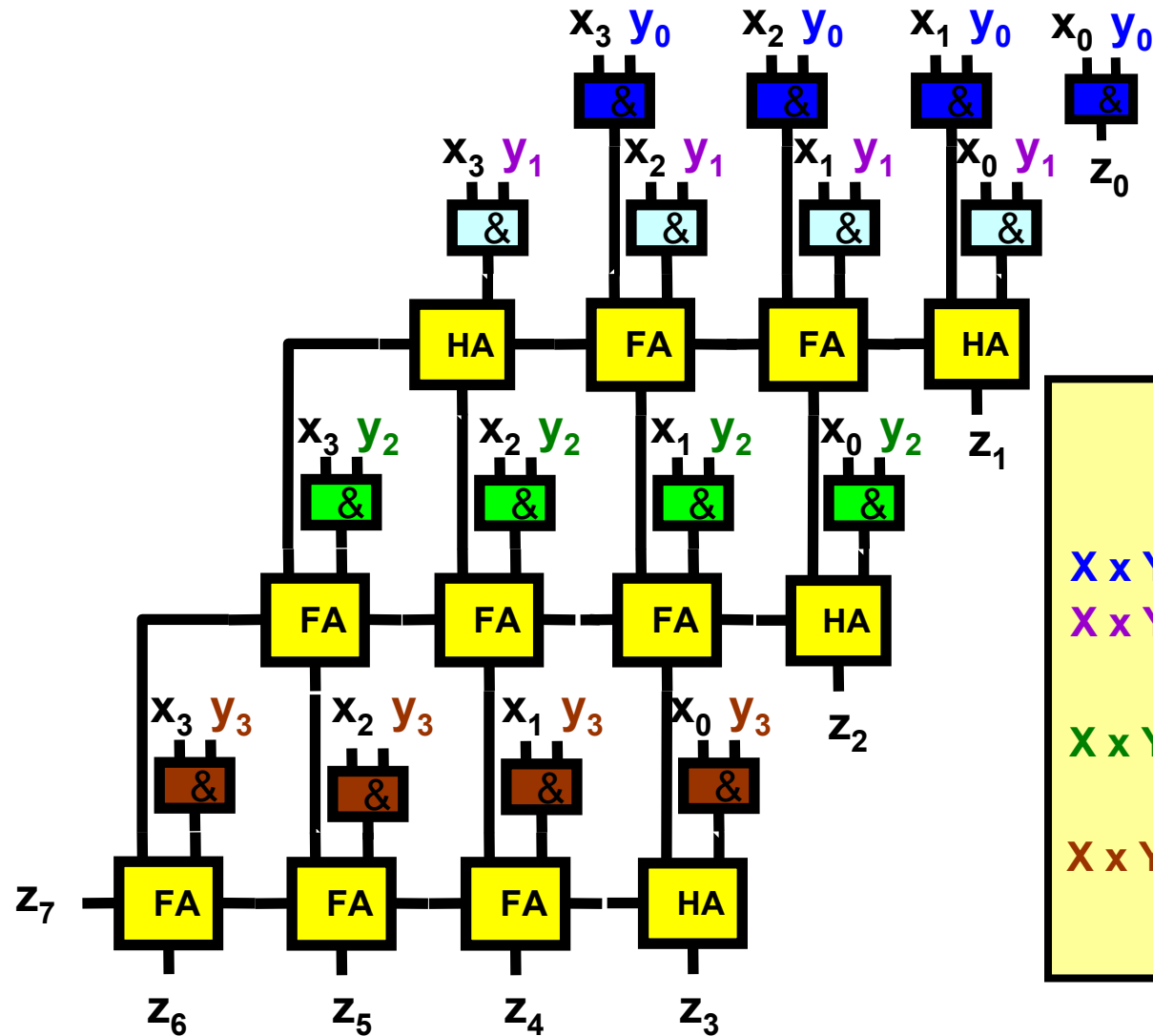
Each partial product formed by bitwise AND operation

Partial products are shifted before being added

■ **Conclusion:** similar to decimal multiplication

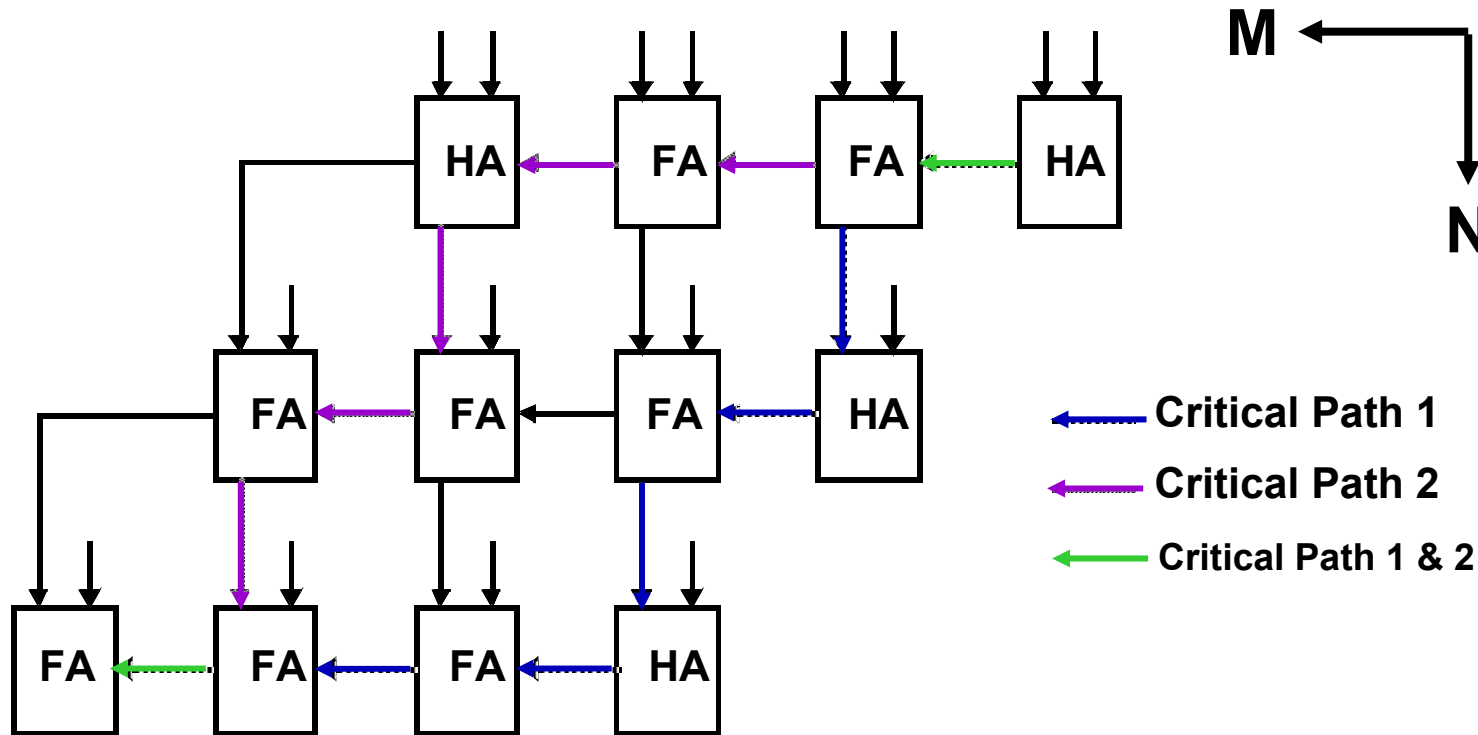
■
$$X_{10} \times Y_{10} = \sum_{i=0}^{M-1} \left(\sum_{j=0}^{N-1} X_i Y_j 10^{i+j} \right)$$

The Array Multiplier



X	1 0 1 0	
Y	1 0 1 1	
	<hr/>	x
$X \times Y_0 \times 2^0$	1 0 1 0	
$X \times Y_1 \times 2^1$	1 0 1 0	
	<hr/>	+
	0 1 1 1 0	
$X \times Y_2 \times 2^2$	0 0 0 0	
	<hr/>	+
	1 0 1 1 1 0	
$X \times Y_3 \times 2^3$	1 0 1 0	
	<hr/>	+
X x Y	1 0 0 1 1 1 0	

The MxN Array Multiplier — Critical Path



- $t_{\text{mult}} \approx [(M - 1) + (N - 2)]t_{\text{carry}} + (N - 1)t_{\text{sum}} + t_{\text{and}}$
- **Requires comparable carry and sum delays**

Adder Cells in Array Multiplier

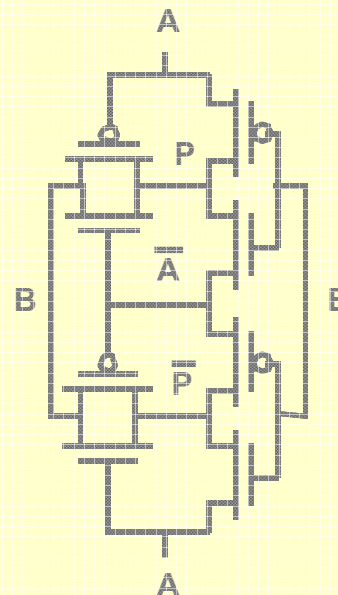
Identical Delays for Carry and Sum

A	B	C_i	S	C_o	Carry status
0	0	0	0	0	delete
0	0	1	1	0	delete
0	1	0	1	0	propagate
0	1	1	0	1	propagate
1	0	0	1	0	propagate
1	0	1	0	1	propagate
1	1	0	0	1	generate
1	1	1	1	1	generate

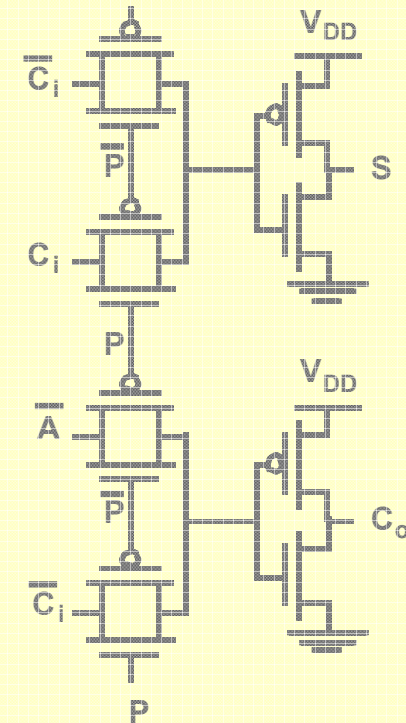
$$P = A \oplus B$$

If $P = 1$ then $S = \bar{C}_i$, $C_o = C_i$

If $P = 0$ then $S = C_i$, $C_o = A$

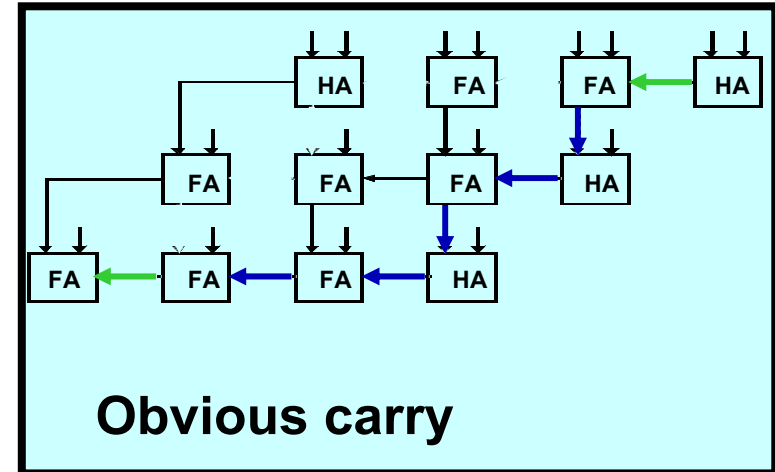
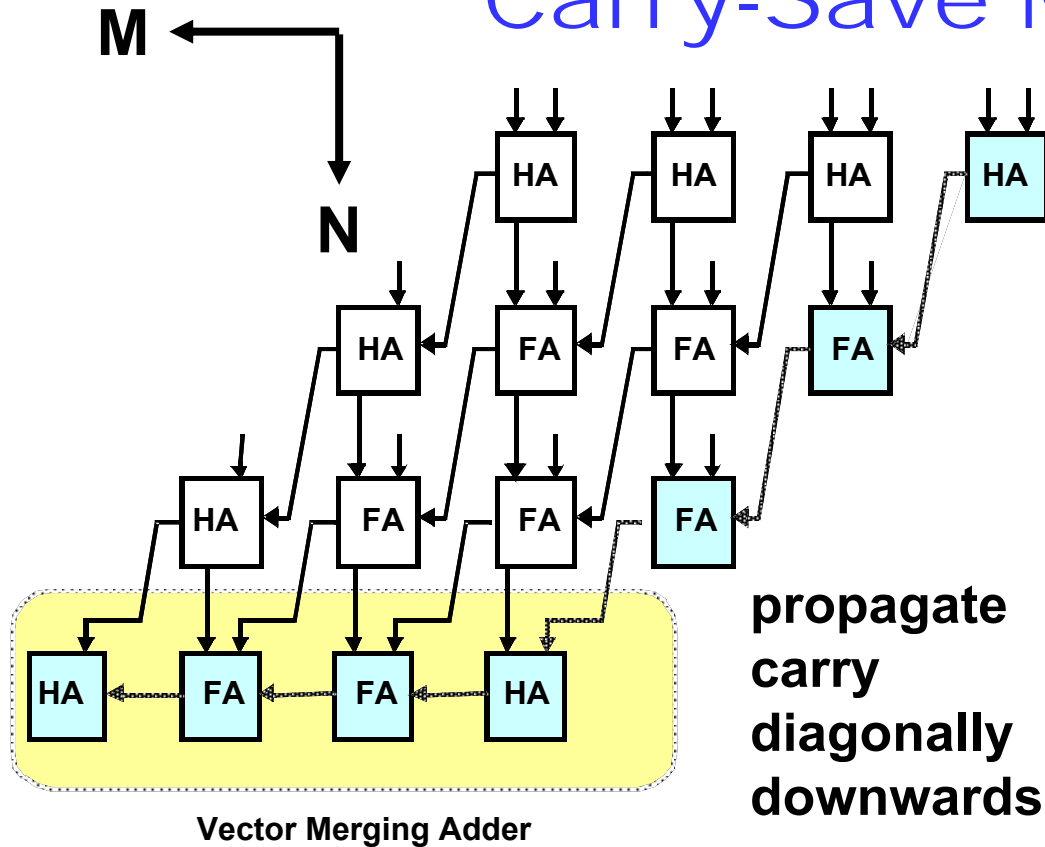


2 x transmission gate XOR for P , \bar{P}
(Fig 11.17)



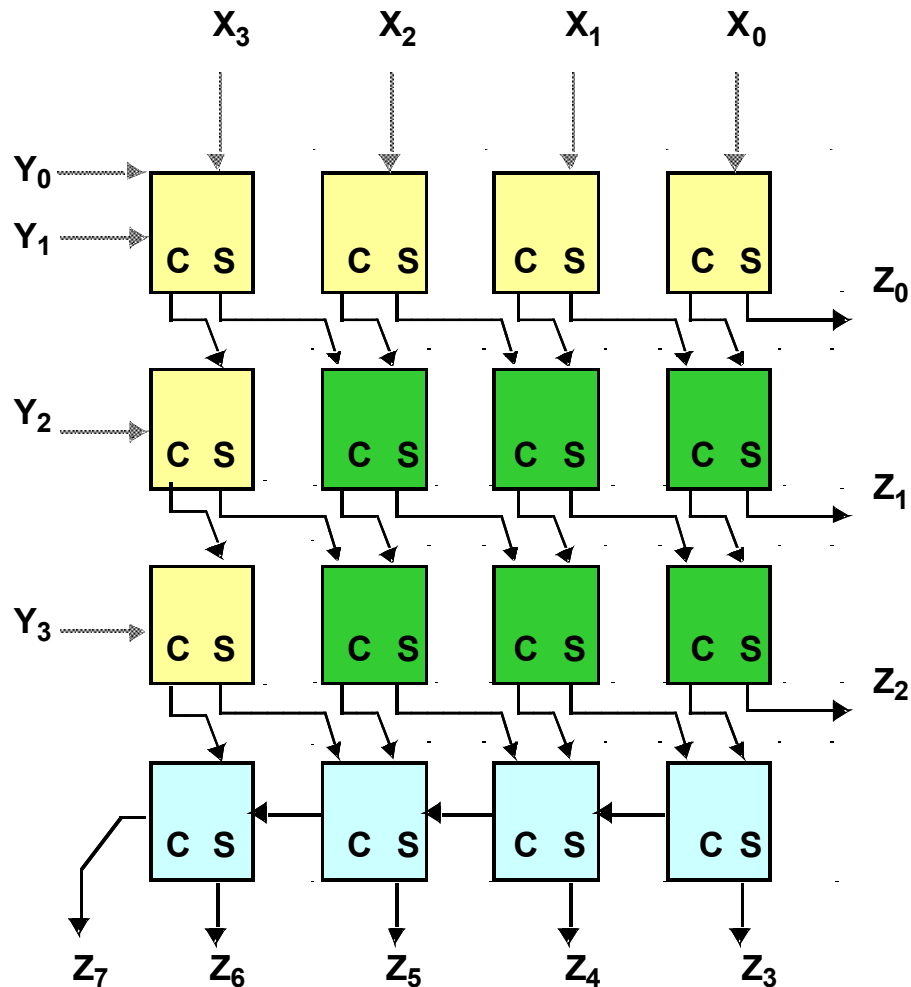
Multiplexers for S , C_o

Carry-Save Multiplier



- $t_{\text{mult}} \approx (N - 1)t_{\text{carry}} + t_{\text{and}} + t_{\text{merge}}$ (assuming $t_{\text{add}} \approx t_{\text{carry}}$)
- Use **fastest possible adder** for final vector merging
- Will be larger, use more power, etc, **but need only one row!**

Multiplier Floorplan



- HA Multiplier Cell
- FA Multiplier Cell
- Vector Merging Cell

X signals are routed vertically across each column (“broadcast”)

Y signal are broadcasted horizontally across rows

Regularity!

Multipliers — Summary.

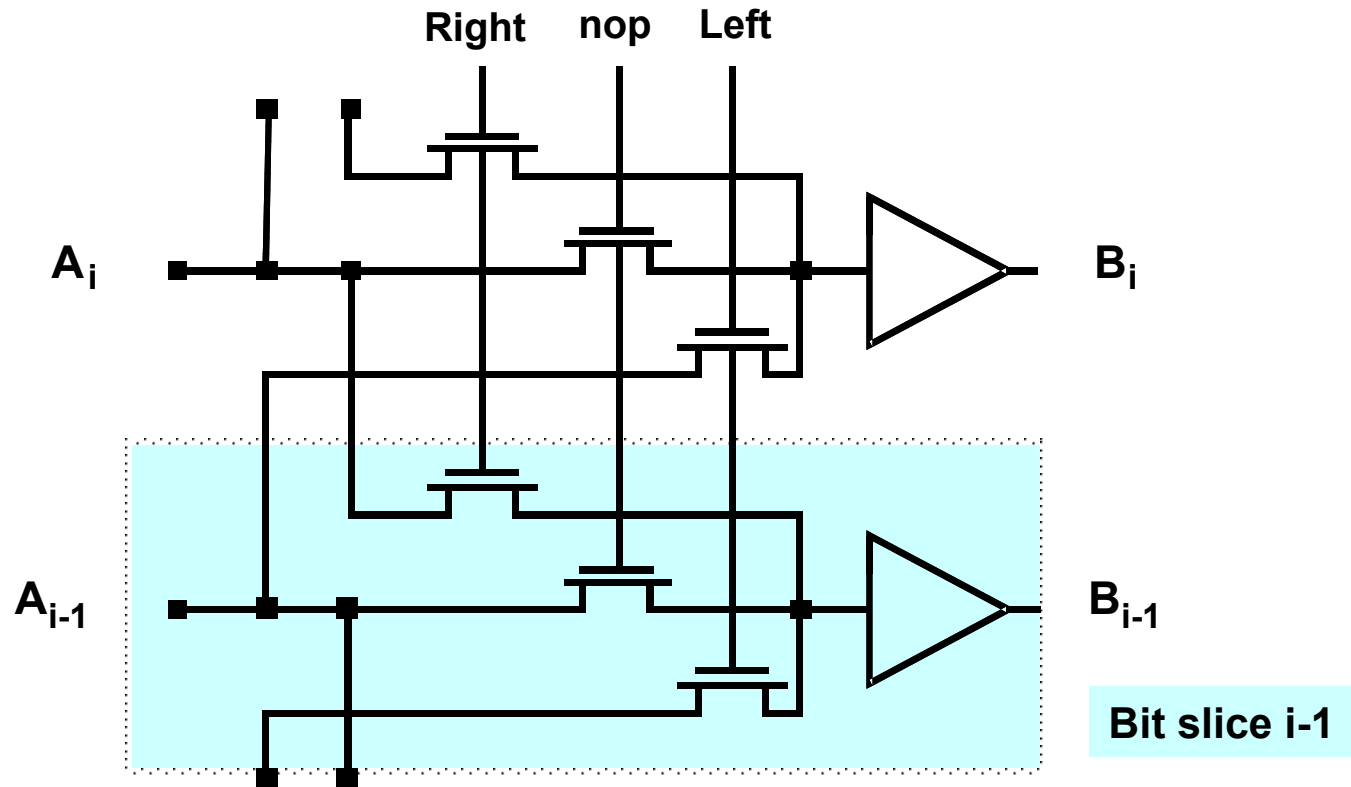
- **Optimization Goals Different Vs Binary Adder**
- **Once Again: Identify Critical Path**
- **Other possible techniques**
 - **Logarithmic versus Linear (Wallace Tree Mult)**
 - **Data encoding (Booth)**
 - **Pipelining**

GLIMPSE AT SYSTEM LEVEL OPTIMIZATION

Shifter Design

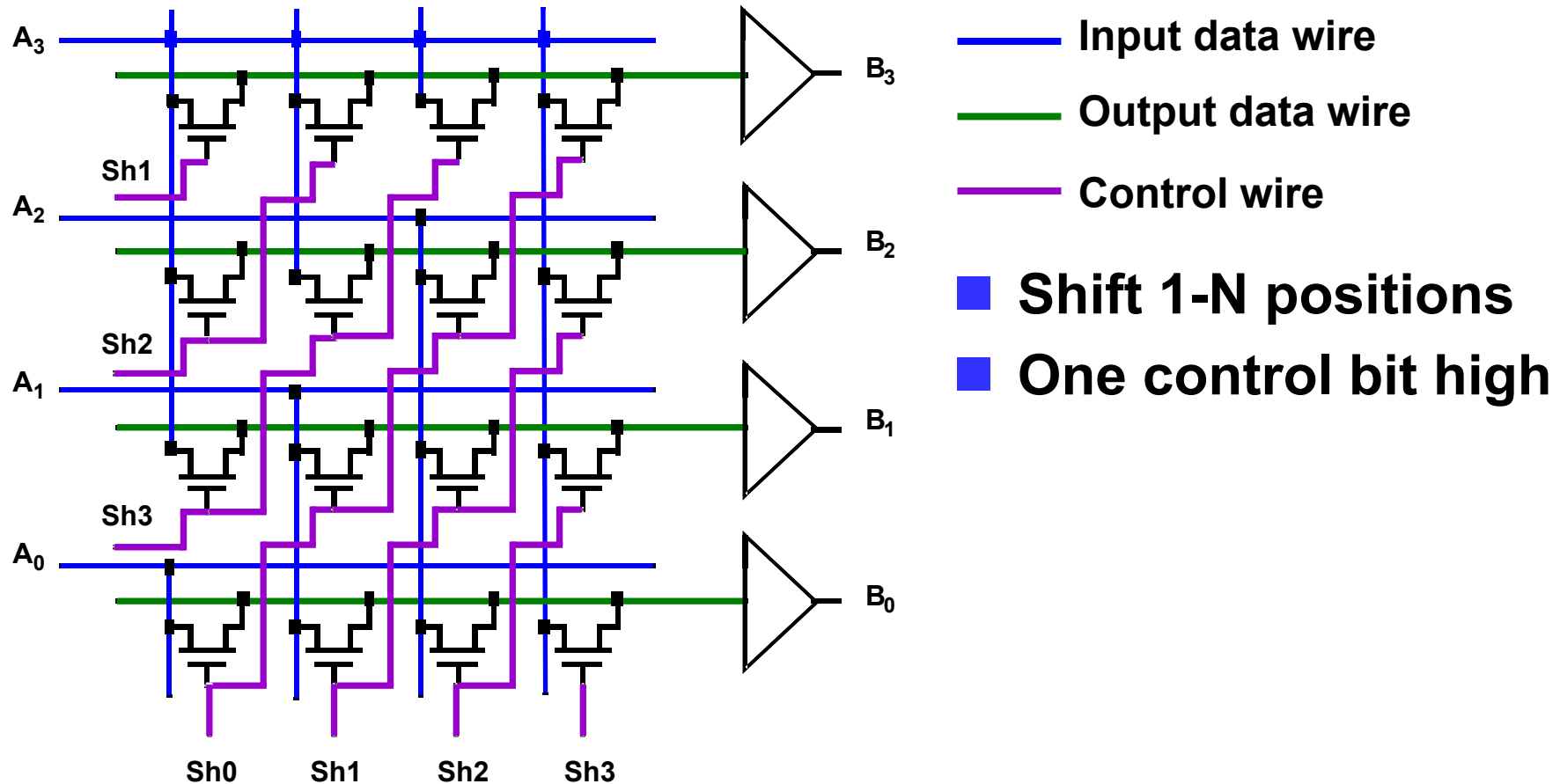
- Shifters are fundamental building blocks **too**
 - Floating point units
 - Scalers
 - Multiplication by constant numbers (add and shift)
 - ...
- Constant shifting is only interconnect
- Programmable shifting requires active circuitry
- Usually dominated by interconnect
- Architectures
 - ✓ ■ Barrel Shifter
 - ✓ ■ Logarithmic Shifter
 - ...
- Design trade-offs, optimization
 - Architecture level, Logic level, Circuit level, Layout level
 - Simpler compared to Adder, Multiplier, hence less rewarding
- Good example of pay-off of structural design

The Binary Shifter



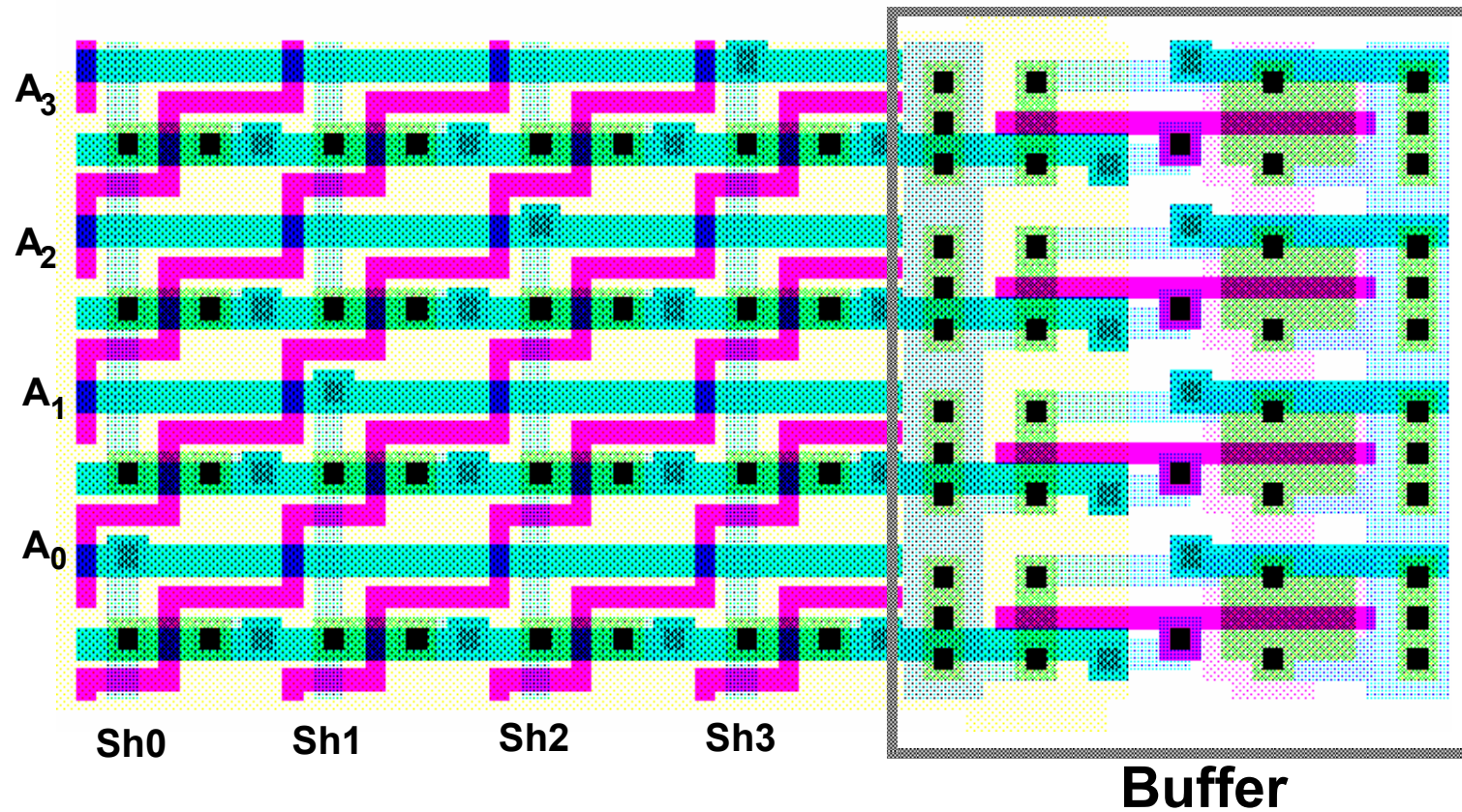
- Multibit shifters by cascading
- M stages for M -bit shift
- Complex and slow for larger M
- More structured approach needed

The Barrel Shifter

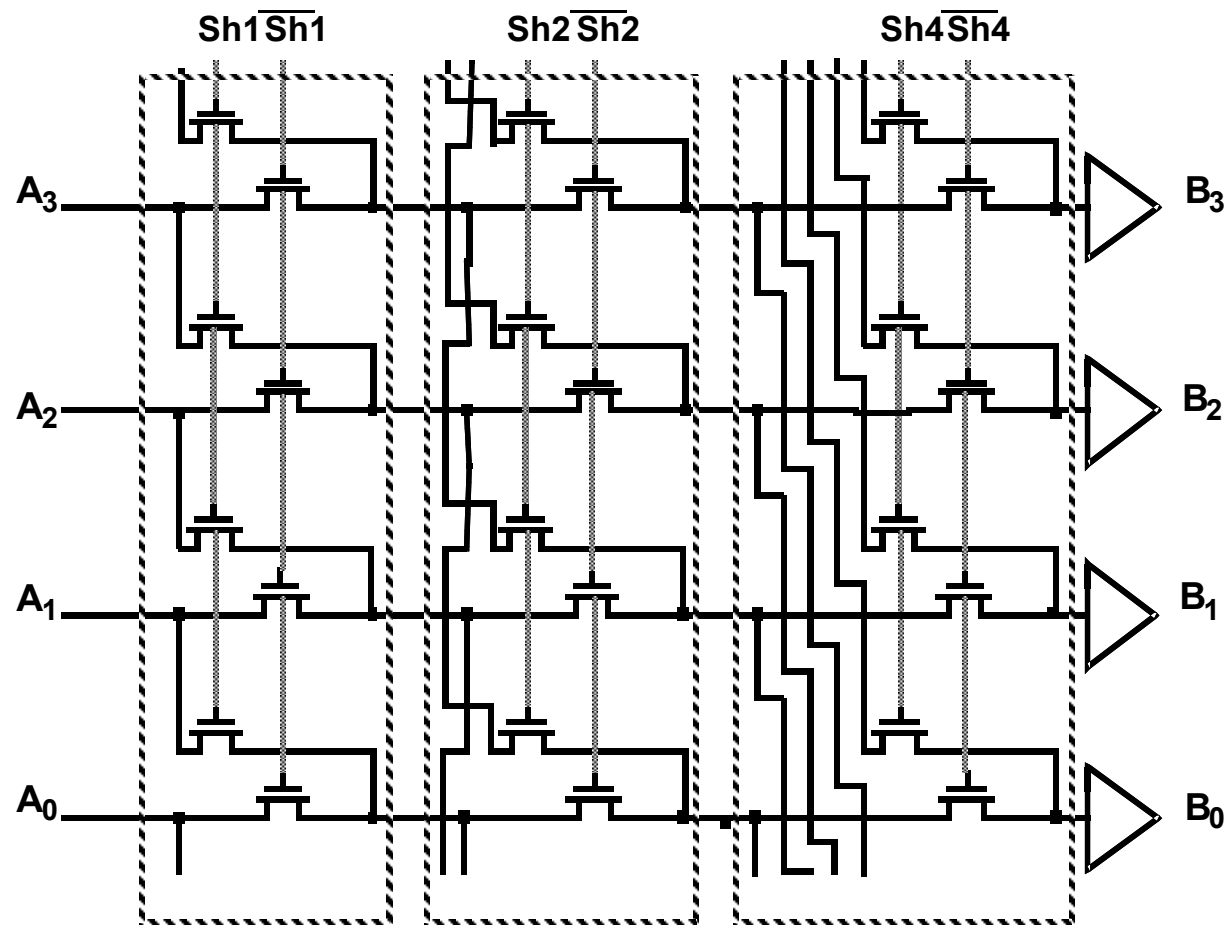


- Need **M stages** for M-bit shift
- Signal passes only one pass-transistor => **delay?**
- **Area** Dominated by Wiring, not (always) by # transistors

4x4 Barrel Shifter

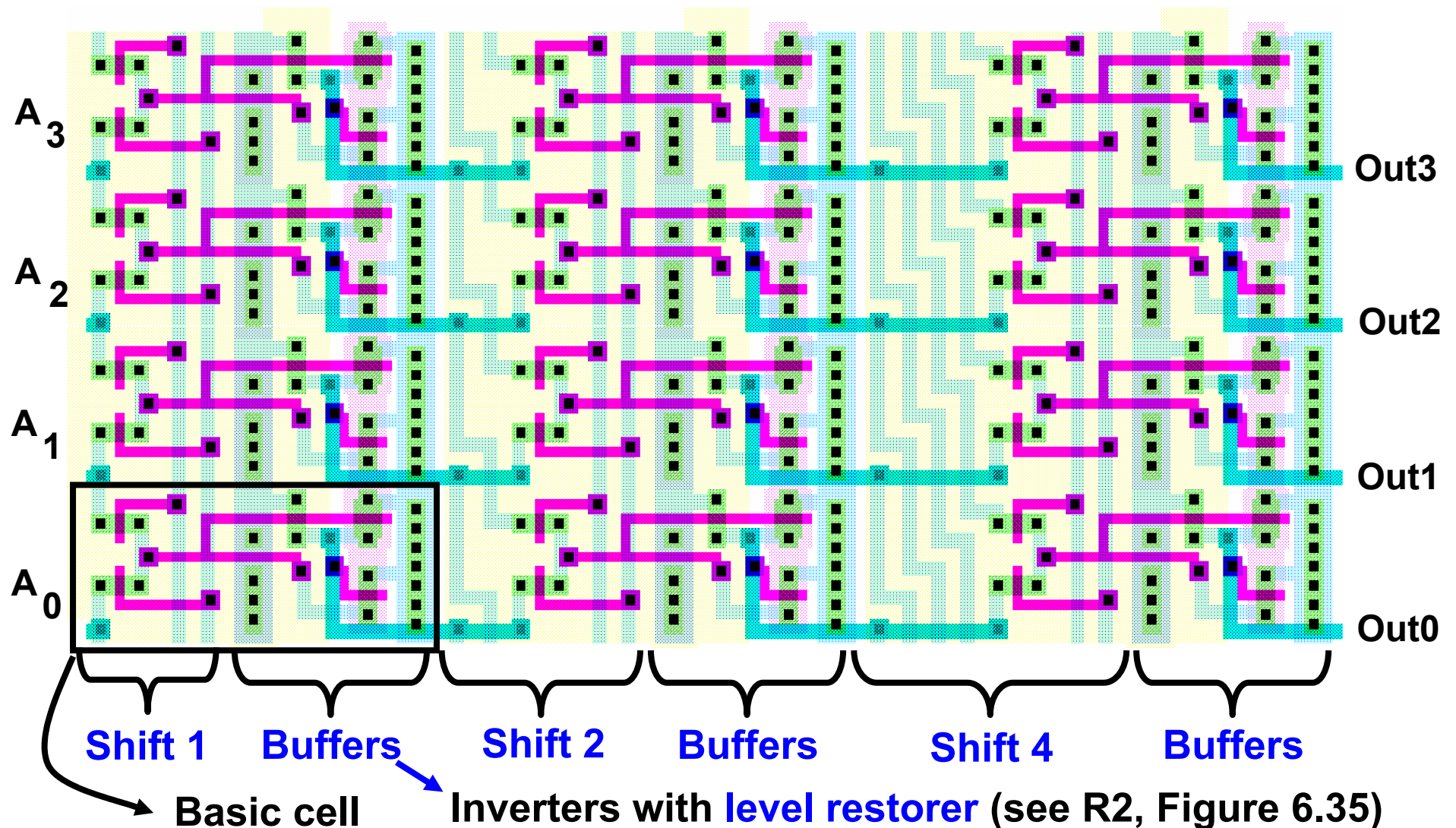


Logarithmic Shifter



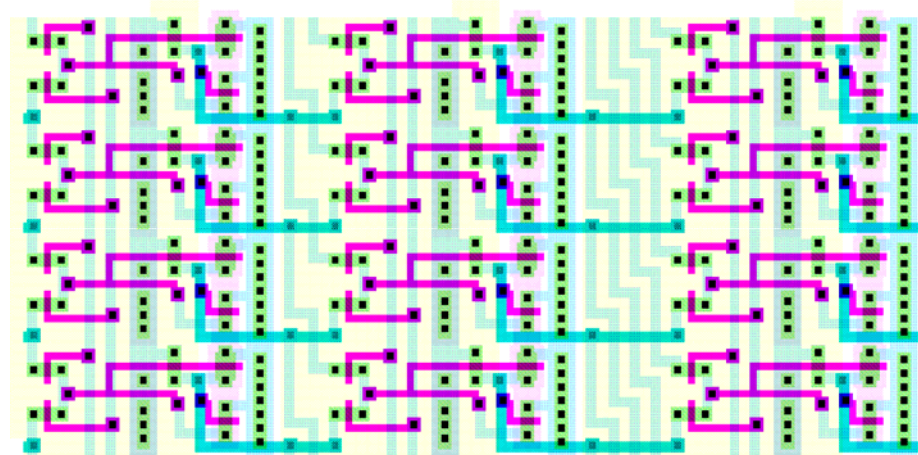
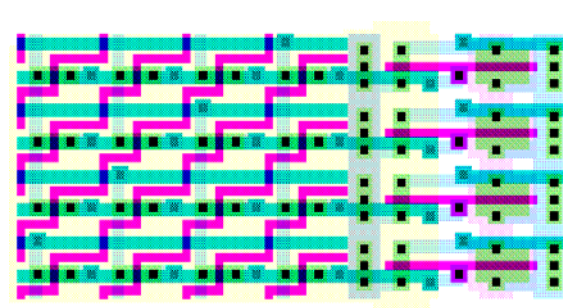
- Section i shifts $2^{(i-1)}$ bits
- Need only $\log_2 M$ stages for M -bit shift

0-7 bit Logarithmic Shifter



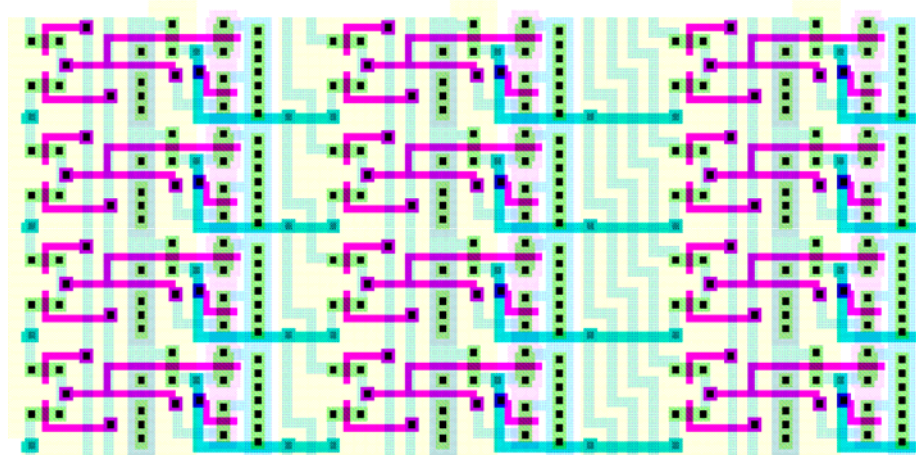
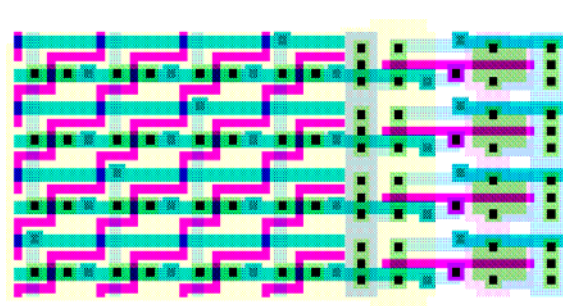
Exercise: decipher layout of basic cell and draw transistor circuit

Size Comparison



- M is maximum bit displacement, $K = \log_2 M$
- For large M , width is dominated by vertical metal wires
- Disregard buffer size, **only count vertical wires**
- Barrel shifter needs 1 control and 1 data wire per stage
- # Wires: $2M$
- Log shifter needs 2 control + 2^{i-1} data wires for stage i
- # Wires: $2K + (1 + 2 + 4 + \dots + 2^{K-1}) = 2K + 2^K - 1 = 2 \log_2 M + M - 1$
- **Log shifter will have smaller area for larger M !**

Speed Comparison

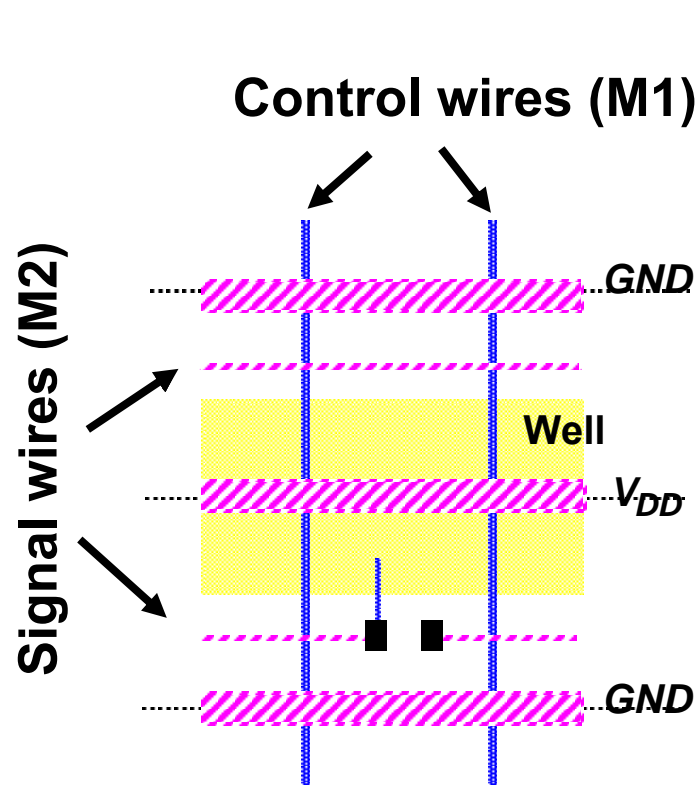


Exercise

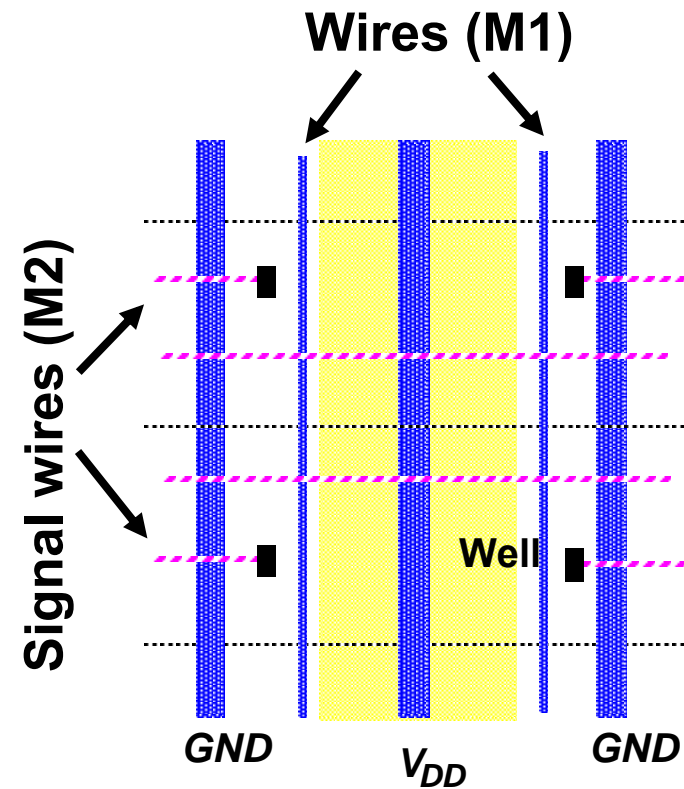
- Discuss the relative speeds of both shifters, as a function of M (see discussion in book). Consider:
 - Number of sections
 - Input capacitance at the buffers (including diffusion areas of the driving pass-transistors)
 - Number of buffers (necessity of buffers)
 - The number of pass-transistors the signal has to pass
 - ...

Layout Strategies (regularity)

Layout Strategies for Bit-Sliced Datapaths

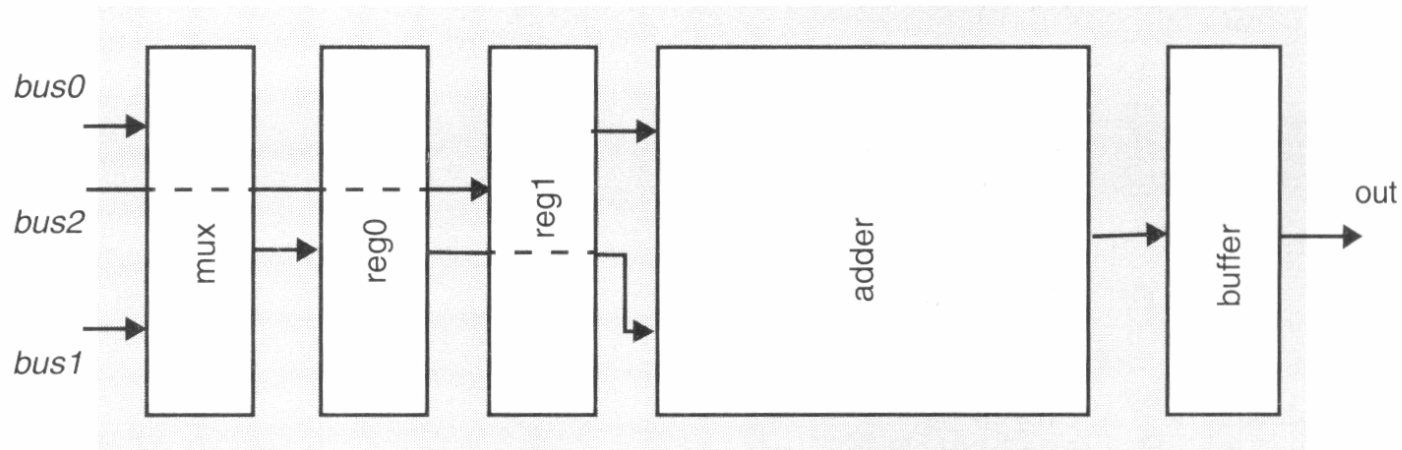


Approach I —
Signal and power lines parallel

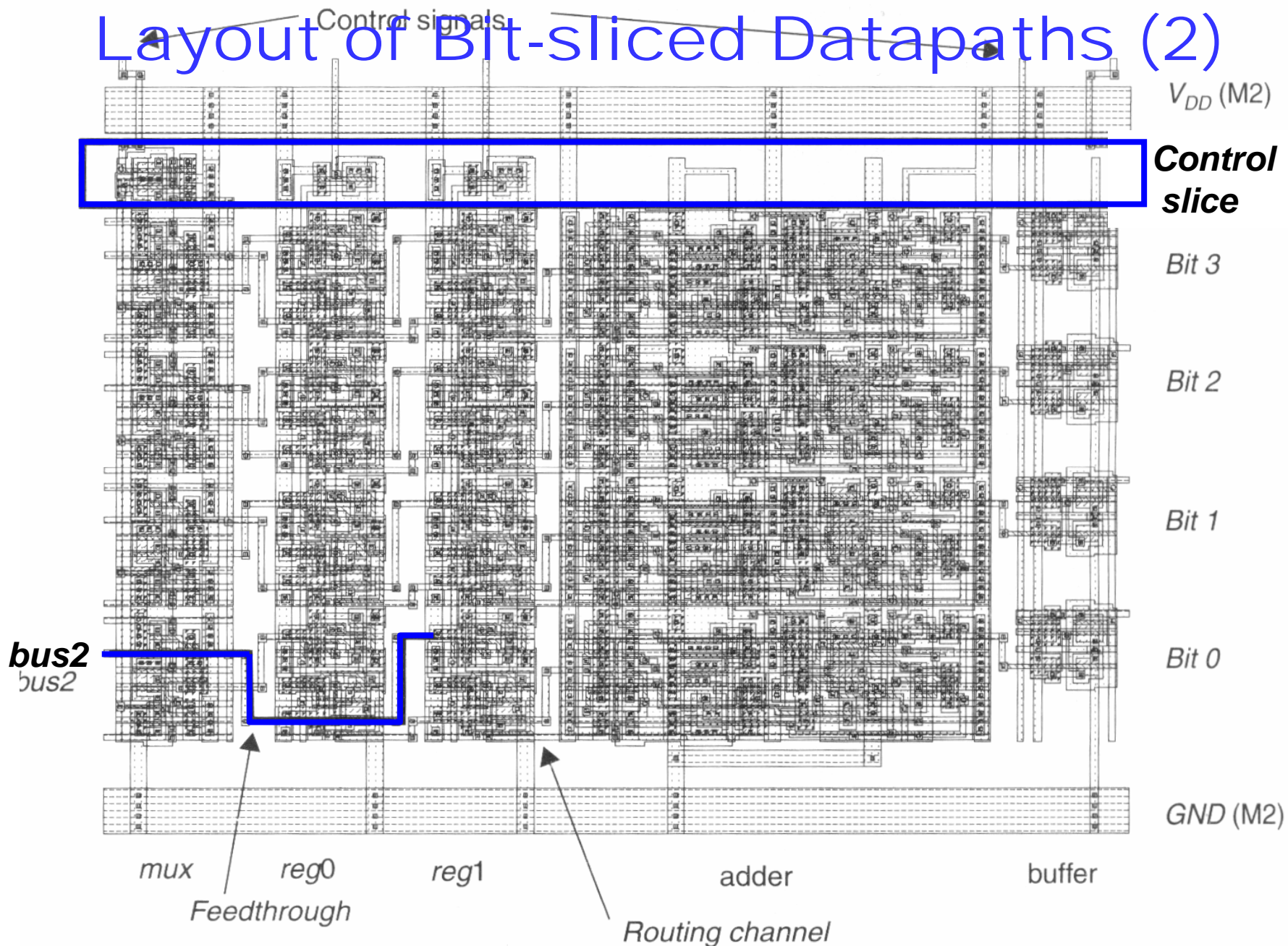


Approach II —
Signal and power lines perpendicular

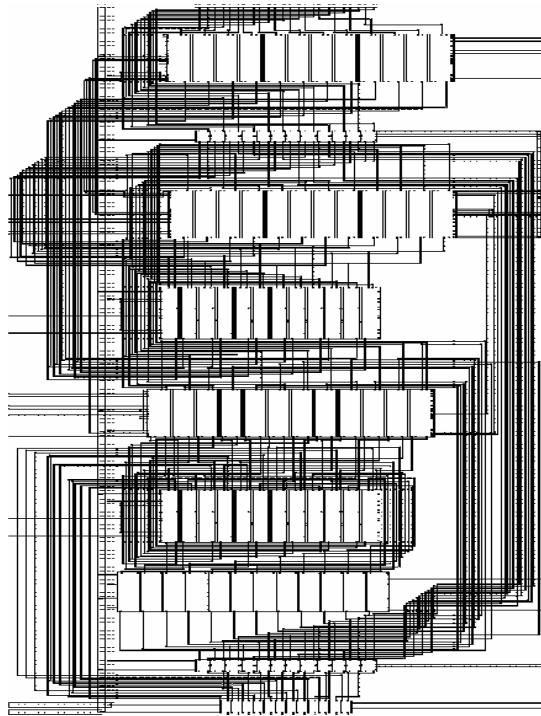
Layout of Bit-sliced Datapaths



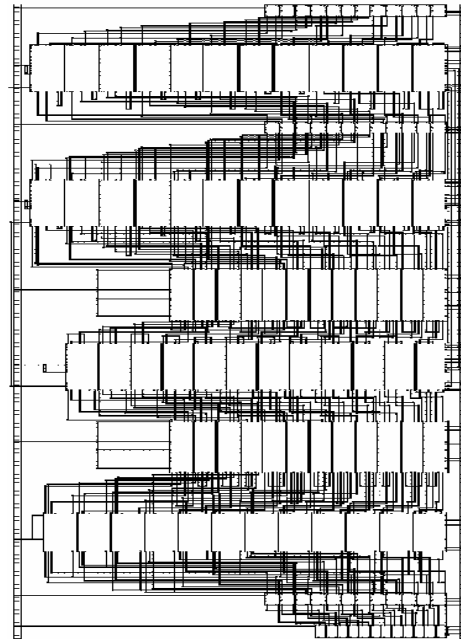
Layout of Bit-sliced Datapaths (2)



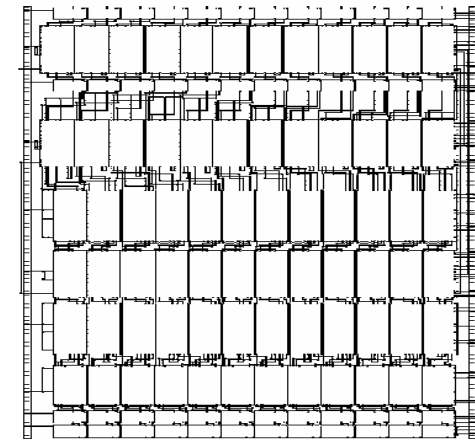
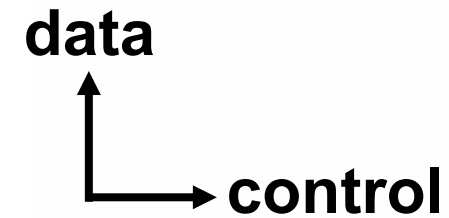
Layout of Bit-sliced Datapaths (3)



Unoptimized
Area: 4.2mm²



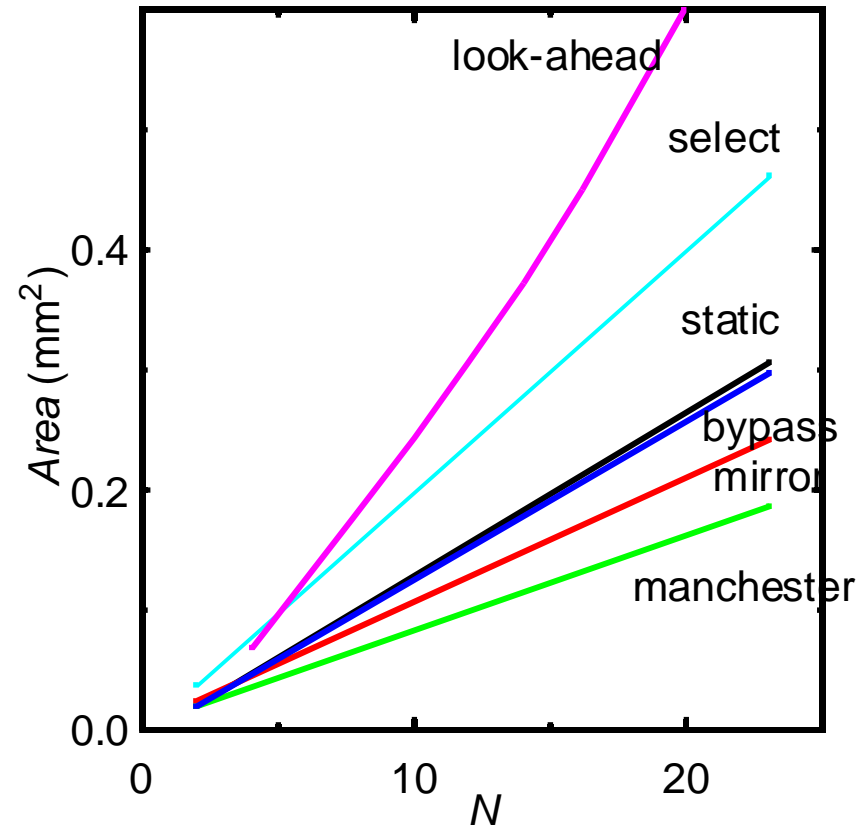
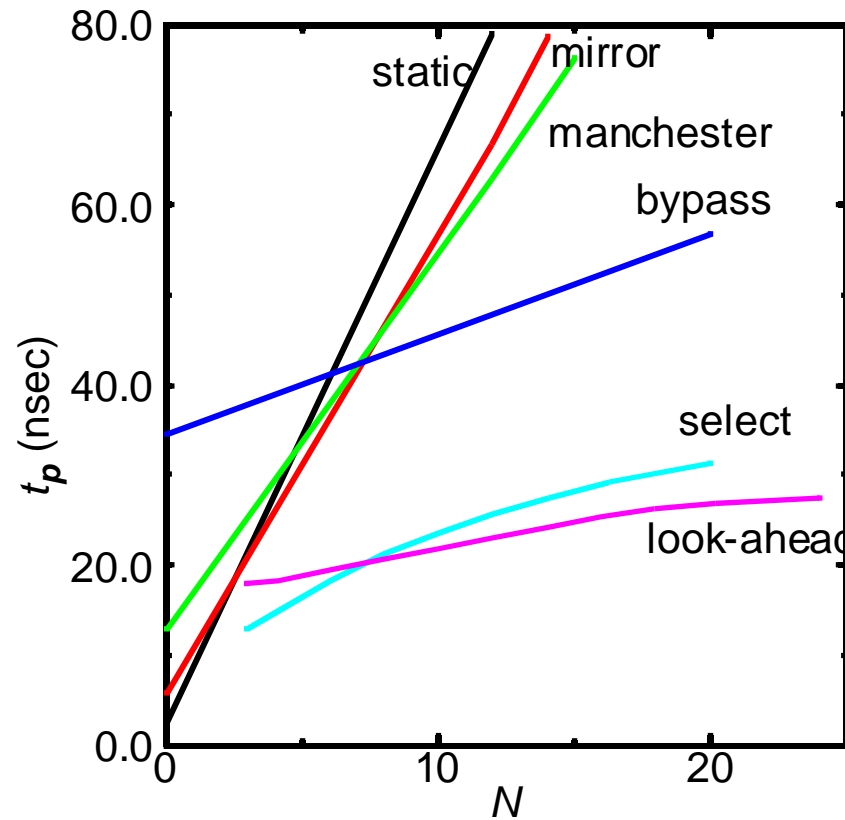
With feedthroughs
Area: 3.2mm²



+ Equalized cell height
Area: 2.2mm²

- **Good layout really counts!**
- **Feedthroughs less (but still) useful with multiple metal layers**

Design as a Trade-Off



VLSI Design

- Select **right structure**
- Determine and optimize **critical timing path** for speed
- Optimize rest for **area** (cost) and/or **power** and/or **design time**
- Consider **layout** aspects

Regularity and **modularity** are a VLSI designer's best friends

Summary

- **Background on Modular Design**
 - **Hierarchy, reuse, regularity**
 - **Architecture, bit-slicing**
- **Adder Design**
- **Multiplier Design**
- **Shifter Design**
- **Layout Strategies (regularity)**
- **Design as a Trade-Off**

**Got further appreciation of some
system level design issues?**

The End

